Introduction

CSE 132 Introduction to Computer Engineering

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Instructional Staff

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- Office hours: TBD (see web page)
- Appointments for Roger: contact via Piazza

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Course Web Page

- cse132.engineering.wustl.edu
- It is a work in progress ...
 - Does contain calendar
 - Will contain studios and assignments
- Documents grading, collaboration, and late policies
- Contains documentation on languages (Java, C) and tools (Eclipse, Git, Arduino IDE)

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What is this class about?

- Organization will be like CSE 131
 - 1.5 hrs/wk lecture (on Wednesdays)
 - 1.5 hrs/wk studio (on Mondays)
- The material includes
 - Basic computer capabilities (I/O, esp. custom I/O)
 - Demystifying how computer systems operate
 - More than one machine, more than one type of machine
 - Design decisions that include both software and hardware

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Some High-level Goals for CSE 132

- Introduce CoE concepts (so those who should be CoE students know what that is)
 - Do this while ensuring relevance to CS students
- Introduce the concept that not all computers are desktop/laptop class machines
 - Computing happens in many different form factors
 - Vehicle for 132 will be an 8-bit microcontroller + standard desktop environment (Java/Eclipse from CSE 131)
- Introduce distributed concurrency (more than one thing going on at a time)
- Recurring theme throughout semester will be the representation of information

Typical Module Sequence

- Lecture
 - Here in Hillman

be staffed by TAs

- Studio

 In Urbauer labs (attendance is required!)
- Assignment
 Demos in office hours or beginning of studio
- Help

 A number of help sessions will get scheduled and
 - Piazza (all the TAs have instructor access)

Demos in Lab

- Lab time on Monday is primarily for studio — Therefore, we reserve most of that time
- Demos
 - Must be in the first 15 min. of lab time (by 1:15)
 - You must be ready to demo when you walk in
 - Help will be available ahead of class, but not in lab
 - No "re-demo" options
 - Have your assignment ready to demo when you request a TA to check it out

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Two Compute Platforms

- Java on laptop or lab machines, using Eclipse as the development environment (just like 131)
- ``C" on Arduino machine
 - Actually a subset of C, and subset is very close to the Java you are familiar with
 - Physical computer is 8-bit machine running at only 16 MHz (over 100 times slower than desktop PC)
 - 16 Kbytes of program memory
 - 2 Kbytes of data memory
 - No keyboard or display
 - Wonderful community of users, doing lots and lots!

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Arduino Programs

- Community calls them "sketches"
- Composed of the basic structure below void setup() {

// insert startup code here, will execute once

void loop() {

}

}

// insert main code here, will execute over and over

Arduino Timing

- Argument is integer number of milliseconds

- Returns the number of milliseconds since last

Later in semester we will use micros()

- Return type is `long int`, which is 32 bits or 4 bytes

Returns number of microseconds since last reset

Use delay() library routine

Use millis() library routine

reset of Arduino

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Hello World • First complete Arduino program void setup() { Serial.begin(9600); //startup comm. link to PC Serial.println("Hello world!"); } void loop() {

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}

Arduino Printing Printing goes to Serial Monitor in Arduino IDE Serial.begin(9600) in setup() initializes port and sets baud rate (communication speed) How do we print? Use Serial.print() and Serial.println() Argument can be any type Serial.println("String to print"); Serial.print(14); // no newline included NOTE: cannot do this – Serial.println("X = " + x); because string concatenation is not supported Do this instead – Serial.print("X = "); Serial.print(x);

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Quiz Time

- Go to Canvas and answer the single question for Quiz 1A
- True or False:

Work cannot be "redemoed" and you can't stop a demo session once is has begun. For example, if you find an error in your work during the demo, you will not be allowed to fix the error and have your grade based on the "fixed" work.

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Let's Get Started

- Computational Abstraction
 Finite-state machines
- Information Representation

 In the digital world, this means binary

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Implementing FSMs

- Bubble diagram is just specification
 - Abstract
 - Independent of implementation
- Possible implementations
 - Directly in hardware (take CSE/ESE 260M to learn)
 - Via software (this is what we will do)
- Important considerations
 - How to represent state (we'll use a single variable)
 - What triggers state transitions (many options)

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Summarizing FSMs

- Useful abstraction of computation
 - Says what needs to be done, not how to do it
 - Independent of implementation
 - Easier to reason about
 - Diagrams easier to read/edit/understand than implementation (no matter how you implement)
- We'll use FSMs many times

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An Example a passed microeconomics course b passed macroeconomics course c passed economics survey course d met economics requirement $d = a \cdot b + c$ $a \xrightarrow{b} \xrightarrow{c} d$

Boolean Algebra

- Boolean algebra (named after 19th century mathematician George Boole) lets us manipulate and reason about expressions of propositional calculus
- Systems based on this algebraic theory are called "digital logic systems"
- All modern computer systems fall in this category

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Physical Representation Positive logic convention Binary value (1 or 0) is represented by the voltage on a wire (H or L) true, 1 voltage greater than threshold V_H false, 0 voltage less than threshold V_L Voltage gap between V_H and V_L provides safety margin to limit errors

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Base 2 (binary) works the same way	
$xyz_2 = x \cdot 2^2 + y \cdot 2^1 + z \cdot 2^0$	
$= x \cdot 4 + y \cdot 2 + z$	
x, y, z can each have 2 possible values: 0 or 1 each digit is called a "bit" e.g., 000 0 001 1 010 2 011 3 100 4 101 5 110 6	
111 /	

How do we represent numbers? A positional number system lets us represent integers. E.g., in base 10:

 $xyz_{10} = x \cdot 10^{2} + y \cdot 10^{1} + z \cdot 10^{0}$ $= x \cdot 100 + y \cdot 10 + z$

x, y, z can each have 10 possible values: 0 to 9

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subtract fixed amount (such as midpoint value) generally bad for computation

4-bit Sign-Magnitude

1st bit encodes sign (0 = positive, 1 = negative) bits 2, 3, 4 magnitude \Rightarrow range 0 to 7 (000 to 111)

overall range -7 to +7 what about 1000? -0!

with n bits, use n-1 bits for magnitude range -(2^{n-1} - 1) to +(2^{n-1} - 1)

issues:

two representations for "0", +0 and -0

need significant hardware to support add, subtract

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2's (radix) complement Use negative weight for 1st bit:

 $wxyz = w \cdot (2)^{3} + x \cdot 2^{2} + y \cdot 2^{1} + z \cdot 2^{0}$ $= w \cdot (8) + x \cdot 4 + y \cdot 2 + z$ • overall range -8 to +7 • 1st bit is still sign bit, with 0 = positive and 1 = negative • only one zero: 0000

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Properties of 2's complement

- least significant n-1 bits have unaltered meaning (i.e., standard positional notation and weights apply)
- most significant bit has weight negated (instead of weight 2ⁿ⁻¹, it is weight -2ⁿ⁻¹)
- range -(2ⁿ⁻¹) to +(2ⁿ⁻¹-1)
- negation operation: flip all bits, add 1, throw away carry
- addition/subtraction function normally

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Make binary more human friendly

- Hexadecimal representation base 16
- Commonly called "hex" but don't be confused, it is not base 6, it is base 16
- Character set 0-9, a-f (alternately A-F) – a=10, b=11, c=12, d=13, e=14, and f=15
- C notation is to prefix hex with symbol 0x (e.g., 0x12, 0xa3)

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Positional notation applies

 $xyz_{16} = x \cdot 16^2 + y \cdot 16^1 + z \cdot 16^0$ = x \cdot 256 + y \cdot 16 + z

So $02c_{16} = 0 \cdot (256) + 2 \cdot (16) + 12 = 44_{10}$

or 0x02c, which is the shorthand I will typically use in class

Benefits of Hex

- Real beauty of hex notation is ease with which one can move back and forth between hex and binary, since 16 = 2⁴
- To transform hex number (e.g., 0x3d50) to binary we expand each hex digit to 4 bits of binary:

3 d 5 0 0011 1101 0101 0000

Binary to Hex Transformation

• To transform binary number (e.g., 1001000) to hex we group into 4-bit groups (starting from right) and rewrite each group in hex

100 1000

4 8 = 0x48

• Or, e.g., 110101110

1 1010 1110

1 a e = 0x1ae

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